HDMN: Hybrid Decentralized Military Network

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Abstract: -- Disruption tolerant network technologies are becoming successful solutions that allow wireless devices carried by soldiers to communicate with each other and access the confidential information or command reliably by exploiting external storage nodes. Some of the most challenging issues in this scenario are the enforcement of authorization policies and the policies update for secure data retrieval. Cipher text attribute with encryption is a cryptographic solution to the access issues. However, the problem of applying in decentralized DTNs apply several security and privacy challenges with regard to the attribute revocation, key escrow, and coordination of attributes issued from different authorities., To propose a secure data retrieval scheme using IDEA for decentralized DTNs where many key authorities manage their attributes independently. We demonstrate how to apply the proposed mechanism to securely and efficiently manage the confidential data distributed in the disruption-tolerant military network.

Keywords— Access Attribute-based encryption (ABE), 3DES algorithm, MD5 algorithm, Disruption-tolerant network (DTN), Multiauthority, and Censored.

I. INTRODUCTION

Mobile nodes in military environments such as a battlefield or a hostile region are likely to suffer from intermittent network connectivity and frequent partitions. Disruption-tolerant network (DTN) technologies are becoming successful solutions that allow wireless devices carried by soldiers to communicate with each other and access the confidential information or command reliably by exploiting external storage nodes. Some of the most challenging issues in this scenario are the enforcement of authorization policies and the policies update for secure data retrieval. Cipher text-policy attribute-based encryption is a promising cryptographic solution to the access control issues. However, the problem of applying in decentralized DTNs introduces several security and privacy challenges with regard to the attribute revocation, key escrow, and coordination of attributes issued from different authorities. We propose a secure data retrieval scheme using for decentralized DTNs where multiple key authorities manage their attributes independently. We demonstrate how to apply the proposed mechanism to securely and efficiently manage the confidential data distributed in the disruption-tolerant military network. Military applications require increased protection of confidential data including access control methods. In many cases, it is desirable to provide differentiated access services such that Data access policies are defined over user attributes or roles, which are managed by the key authorities.

ll. LITERATURE SERVERY

Mobile nodes in military environments such as a battlefield or a hostile region are likely to suffer from intermittent network connectivity and frequent partitions. Disruption-tolerant network (DTN) technologies are becoming successful solutions that allow wireless devices carried by soldiers to communicate with each other and access the confidential information or command reliably by exploiting external storage nodes. Some of the most challenging issues in this scenario are the enforcement of authorization policies and the policies update for secure data retrieval. Ciphertext-policy attribute-based encryption() is a promising cryptographic solution to the access control issues. However, the problem of applying in decentralized DTNs introduces several security and privacy challenges with regard to the attribute revocation, key escrow, and coordination of attributes issued from different authorities. Border Surveillance: A dynamic deployment scheme for WSN-based solutions Ramzi Bellazreg1, Noureddine Boudriga1, Khalifa Trimèche 2 and Sunshin An31 University of Carthage Tunisia, 2 Faculty of Science of Tunis Tunisia and 3Korea University ©2013 IEEE Wireless Sensor Networks (WSNs) are based on

elementary sensors that detect the occurrence of particular events in a monitored area. Among the recent critical WSN applications one can find the border surveillance applications. The first aim of this class of applications is to monitor a country border and detect the presence of intruders near the border line. In this paper, we investigate theoretically the effects of natural factors on dynamic deployment scheme of a hierarchical WSN-based solution providing two lines of surveillance. Parameters such as the wind effect, the altitude and velocity of the airplane from which the sensors are thrown are put into equation to optimize the area coverage and WSN connectivity. Then, we propose mathematical models that evaluate the quality of connectivity and coverage of the deployed network and allow planning and dimensioning of a border solution.

Barrier Coverage with Airdropped Wireless Sensors Anwar Saipulla Benyuan Liu Jie Wang Department of Computer Science University of Massachusetts Lowell , MA 01854 USA 2008 IEEE. Barrier coverage of a wireless sensor network aims at detecting intruders crossing the network. It provides a viable alternative for monitoring boundaries of battlefields, country borders, coastal lines, and perimeters of critical infrastructures. Early studies on barrier coverage typically assume that sensors are deployed uniformly at random in a large area. This assumption, while theoretically interesting, may be unrealistic in real applications. We take a more realistic approach in this paper. In particular, we consider that sensors are airdropped from an aircraft along its flying route. We note that wind, geographic terrain, and other factors may cause a sensor to land in a location deviating from its targeted landing point with a random offset. Thus, it is more realistic to assume that sensor nodes are distributed with a normal offset along the deployment line.

lll. PROPOSED SYSTEM

3.1. Key Authorities: They are key generation centers that generate public/secret parameters for. The key authorities consist of a central authority and multiple local authorities. We assume that there are secure and reliable communication channels between a central authority and each local authority during the initial key setup and generation phase. Each local authority manages different attributes and issues corresponding attribute keys to users. They grant differential access rights to individual users based on the users' attributes. The key authorities are assumed to be honest-but-curious. That is, they will honestly execute the assigned tasks in the system; however they would like to learn information of encrypted contents as much as possible.

3.2. Storage node: This is an entity that stores data from senders and provide corresponding access to users. It may be mobile or static. Similar to the previous schemes, we also assume the storage node to be semi trusted, that is honest-but-curious.

3.3. Sender: This is an entity who owns confidential messages or data and wishes to store them into the external data storage node for ease of sharing or for reliable delivery to users in the extreme networking environments. A sender is responsible for defining (attribute based) access policy and enforcing it on its own data by encrypting the data under the policy before storing it to the storage node.

3.5. User: This is a mobile node who wants to access the data stored at the storage node (e.g., a soldier). If a user possesses a set of attributes satisfying the access policy of the encrypted data defined by the sender, and is not revoked in any of the attributes, then he will be able to decrypt the 3DES algorithm and obtain the data. Since the key authorities are semi-trusted, they should be deterred from accessing plaintext of the data in the storage node; they should be still able to issue secret keys to users. In order to realize this somewhat contradictory requirement, the entral authority and the local authorities engage in the arithmetic with master secret keys of their own and issue independent key components to users during the key issuing phase. It prevents them from knowing each other's master secrets so that none of them can generate the whole set of secret keys of users individually.
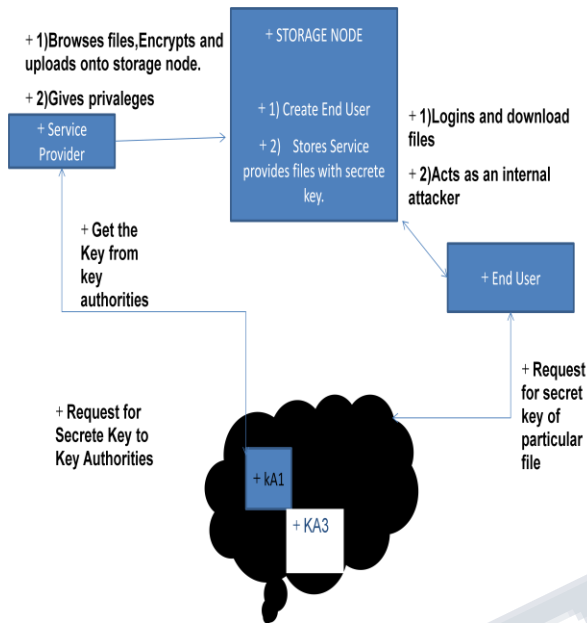
Fig-3.1: Encryption and key generation

Pervious Algorithm CPABE

We propose an attribute-based secure data retrieval scheme using CP-ABE for decentralized DTNs. The proposed scheme features the following achievements. First, immediate attribute revocation enhances backward/forward secrecy of confidential data by reducing the windows of vulnerability. Second, encryptors can define a fine-grained access policy using any monotone access structure under attributes issued from any chosen set of authorities. Third, the key escrow problem is resolved by an escrow-free key issuing protocol that exploits the characteristic of the decentralized DTN architecture. The key issuing protocol generates and issues user secret keys by performing a secure two-party computation (2PC) protocol among the key authorities with their own master secrets. The 2PC protocol deters the key authorities from obtaining any master secret information of each other such that none of them could generate the whole set of user keys alone. Thus, users are not required to fully trust the authorities in order to protect their data to be shared. The data confidentiality and privacy can be cryptographically enforced against any curious key authorities or data storage nodes in the proposed scheme.

DES with MD5 ALGORTHIM

3DES encrypts a 64-bit block of plaintext to 64-bit block of ciphertext. It uses 128-bit key. The algorithm consists of eight identical rounds and a "half" round final Transformation. There are 216 possible 16-bit blocks each operation with the set of possible 16-bit blocks is an algebraic group. Bitwise XOR is bitwise addition modulo 2, and addition modulo 216 is the usual group operation. Some spin must be put on the elements – the 16-bit blocks – to make sense of multiplication modulo 216 + 1, however. 0 is not an element of the multiplicative group.

Confidentiality: In order to protect sensed data and communication exchanges between sensor nodes it is important to guarantee the secrecy of messages. In the sensor network case this is usually achieved by the use of symmetric cryptography as asymmetric or public key cryptography in general is considered too expensive. However, while encryption protects against outside attacks, it does not protect against inside attacks/node compromises, as an attacker can use recovered cryptographic key material to successfully eavesdrop, impersonate or participate in the secret communications of the network. Furthermore, while confidentiality guarantees the security of communications inside the network it does not prevent the misuse of information reaching the base station. Hence, confidentiality must also be coupled with the right control policies so that only authorized users can have access to confidential information.

Integrity and Authentication: authentication is necessary to enable sensor nodes to detect modified, injected, or replayed packets. While it is clear that safety-critical applications require authentication, it is still wise to use it even for the rest of applications since otherwise the owner of the sensor network may get the wrong picture of the sensed world thus making inappropriate decisions. However, authentication alone does not solve the problem of node takeovers as compromised nodes can still authenticate themselves to the network. Hence authentication mechanisms should be "collective" and aim at securing the entire network. In particular, the following requirements must be supported by the key management scheme, in order to facilitate data aggregation and dissemination process: First we focused on the establishment of trust

relationship among wireless sensor nodes, and presented a key management protocol for sensor networks. The protocol includes support for establishing four types of key spersenso rnode: individual keys shared with the base station, pair wise keys shared with individual neighboring nodes, cluster keys shared with a set of neighbors, and a group key shared with all the nodes in the network. We showed how the keys can be distributed so that the protocol can support in-network processing and efficient dissemination, while restricting the security impact of a node compromise to the immediate network neighborhood of the compromised node. Applying the protocol makes it really hard for an adversary to disrupt the normal operation of the network.
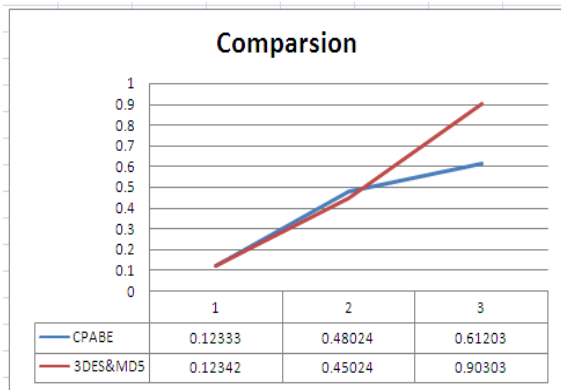
IV. RESULT

Fig 2. Packet Delivery Ratio

	1	2	3
CPABE	0.12333	0.48024	0.61203
3DES&MD5	0.12342	0.45024	0.90303

V. RESEARCH METHOD

Encryption Alolgorithm

1. Take a file packet [N]

2. Encrypt the plaintext blocks using single DES with E key K_1.

3. Now decrypt the output of step 1 DP using single DES with key K_2.

4. Finally, encrypt the output of step 2 using single DES with key K_3.

5. The output of step 3 is the ciphertext. (CT)

6. 128 key by md5 (MD) KD

7. $CT = N K_3 (DP K_2 (E K_1))$

8. $Iblock = CT + KD\backslash$

9. Iblock is send to battalion receiver.

Decryption Algorithm

1. Received $Iblock = CT + KD$

2. First and second secret keys or second and third secret keys are the same.

3. Whichever key.

4. $C = E3 (KD1 (K_1 (N))) = E3 (N)$

5. $C = E3 (KD3 (K_2 (N))) = K_2 (N)$

6. It is possible to use 3DES cipher with a secret 128 bit key.

7. In this case first and third secret keys are the same.

8. $C = K 1(KD2 (K1 (N))$

9. If key match data decryption
$m = D1 (K_2 (KD3(c)))$

VI. METHODOLOGY

We propose an attribute-based secure data retrieval scheme using for decentralized DTNs. The proposed scheme features the following achievements. First, immediate attribute revocation enhances backward/forward secrecy of confidential data by reducing the windows of vulnerability. Second, encryptors can define a fine-grained access policy using any monotone access structure under attributes issued from any chosen set of authorities. Third, the key escrow problem is resolved by an escrow-free key issuing

protocol that exploits the characteristic of the decentralized DTN architecture. The key issuing protocol generates and issues user secret keys by performing a secure two-party computation among the key authorities with their own master secrets. The key authorities from obtaining any master secret information of each other such that none of them could generate the whole set of user keys alone. Thus, users are not required to fully trust the authorities in order to protect their data to be shared. The data confidentiality and privacy can be cryptographically enforced against any curious key authorities or data storage nodes in the proposed scheme.

VII. CONCLUSION

The corresponding attribute group keys are updated and delivered to the valid attribute group members securely (including the user). In addition, all of the components encrypted with a secret key in the ciphertext are reencrypted by the storage node with a Random and the ciphertext components corresponding to the attributes are also reencrypted with the updated attribute group keys. Even if the user has stored the previous ciphertext exchanged before he obtains the attribute keys and the holding attributes satisfy the access policy, he cannot decrypt the pervious ciphertext.

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